

# Matching with Contracts

Hatfield and Milgrom (AER, Sep 2005)

summary by N. Antić

Examines connection between two-sided matching and auction/contract theory. Echenique (2012, AER) shows this model can be embedded in the Kelso-Crawford (1982, ECTA) framework. The extension by Hatfield-Kojima (2010, JET) does not embed in the older framework, thus this is a good workhorse model.

## Basic Model

- Finite sets of doctors and hospitals  $D$  and  $H$ , respectively
- $X$  is a finite set of contracts (generic)
  - A contract  $x \in X$  matches a doctor,  $x_D$ , and hospital,  $x_H$
  - $x$  may contain more information, e.g., doctor's wage
  - $A \subset X$  is an allocation if  $\forall (d, h), (d', h') \in A, d \neq d'$
- $\succ_d$  is a total order on  $X_d \cup \{\emptyset\}$ ,  $X_d := \{x \in X \mid x_D = d\}$ 
  - Contract  $x$  is acceptable to doctor  $d$  if  $x \succ_d \emptyset$
  - Given  $Y \subset X$ , let  $C_d(Y) = \max_{\succ_d}(X_d \cap Y \cup \{\emptyset\})$ 
    - $C_D(Y) = \cup_d C_d(Y)$  denotes accepted contracts
    - $R_D(Y) = Y \setminus C_D(Y)$  denotes rejected contracts
- $\forall Y \subset X$ , hospital preferences are given by a choice correspondence,  $C_h(Y) \subset Y_h := \{x \in Y \mid x_H = h\}$ , where for any  $x, x' \in C_h(Y)$ ,  $x \neq x' \Rightarrow x_D \neq x'_D$ 
  - Let  $C_H(Y) = \cup_h C_h(Y)$ ,  $R_H(Y) = Y \setminus C_H(Y)$
- A set of contracts  $A \subset X$  is a stable allocation if
  - S1.  $C_D(A) = C_H(A) = A$ , and
  - S2.  $\nexists h \in H, Y \subset X$  s.t.  $Y = C_h(A \cup Y) \subset C_D(A \cup Y)$

**Theorem 1.** If  $(X_D, X_H) \subset X^2$  solves the system:

$$\begin{aligned} X_D &= X \setminus R_H(X_H), \\ X_H &= X \setminus R_D(X_D), \end{aligned}$$

then  $A := X_H \cap X_D$  is a stable allocation and  $A = C_H(X_H) = C_D(X_D)$ . Conversely, for any stable  $A$ , there exist  $X_D, X_H$  satisfying the system such that  $A = X_H \cap X_D$ .

*Proof.* ( $\Rightarrow$ )  $A := X_H \cap X_D = X_D \setminus R_D(X_D) = C_D(X_D)$ ; similarly  $A = C_H(X_H)$ . By revealed preference  $A = C_D(X_D) = C_D(A) = C_H(A) = C_H(X_H)$ , thus  $A$  satisfies S1. Fix  $h$  and some  $Y_h \subset C_D(A \cup Y_h)$ . Since  $A = C_D(X_D)$ , by doctors' revealed preferences,  $Y_h \cap X_D \subset C_D(X_D)$  and thus

$$Y_h \cap R_D(X_D) = Y_h \cap X_D \cap R_D(X_D) \subset C_D(X_D) \cap R_D(X_D) = \emptyset.$$

Hence,  $Y_h \subset X \setminus R_D(X_D) = X_H$  and so if  $Y_h \neq C_h(A)$ , then  $Y_h \prec_h C_h(X_h) = C_h(A)$ . Thus  $Y_h \neq C_h(A \cup Y_h)$  and S2 holds.

( $\Leftarrow$ ) Let  $A$  be a stable allocation. Note that  $A = C_H(A) = C_D(A)$ . Define:

$$\begin{aligned} X_H &= \bigcup_{d \in D} \{x \in X_d \cup \{\emptyset\} : x \succeq_d A_d\}, \\ X_D &= \bigcup_{d \in D} \{x \in X_d \cup \{\emptyset\} : x \preceq_d A_d\}. \end{aligned}$$

Clearly  $A, X_D \setminus A, X_H \setminus A$  is a partition of  $X$ . If  $C_h(X_H) \neq A_h$ , then  $Y = C_h(X_H)$  violates S2. Thus  $C_H(X_H) = A$ , since  $C_h(X_H) = A_h$  for all  $h$ . Hence:

$$\begin{aligned} X \setminus R_H(X_H) &= X \setminus (X_H \setminus A) = X_D, \text{ and} \\ X \setminus R_D(X_D) &= X \setminus (X_D \setminus C_D(X_D)) = X_H, \end{aligned}$$

since by definition  $C_D(X_D) = A$ . □

- Elements of  $X$  are substitutes for  $h$  if  $R_h(Y) \subset R_h(Z)$  for any  $Y \subset Z \subset X$

- i.e., if  $R_h$  is isotone (order-perserving) with respect to  $\subset$
- New contracts make hospital weakly less interested in old
- Coincides with demand theory definition in the appropriate setting (Theorem 2)

## Generalized Deferred Acceptance (DA) Algorithm

- Generalized DA algorithm,  $F : X \times X \rightarrow X \times X$

$$F(X_D, X_H) = (X \setminus R_H(X_H), X \setminus R_D(X \setminus R_H(X_H)))$$

- $R_D$  (by revealed preference) and  $R_H$  (by assumption) are isotone, and thus  $F$  is also isotone wrt  $\geq$ , defined as:

$$F(X_D, X_H) \geq F(Y_D, Y_H) \text{ iff } Y_D \subset X_D \text{ and } X_H \subset Y_H$$

**Theorem (Tarski).** If  $(L, \geq)$  is a complete lattice and  $f : L \rightarrow L$  is isotone, then  $f$  has a fixed point. Further, if  $P$  is the set of fixed points of  $f$ , then  $(P, \geq)$  is a complete lattice.

- A special case of this is the existence theorem in the paper

**Theorem 3.** Suppose contracts are substitutes for hospitals. Then:

- (a) The set of fixed points of  $F$  is a complete (finite) lattice
- (b) Starting at  $(X, \emptyset)$ , iteration on  $F$  converges monotonically to the highest fixed point  $(\underline{X}_D, \underline{X}_H)$
- (c) Starting at  $(\emptyset, X)$ , iteration on  $F$  converges monotonically to the lowest fixed point  $(\underline{X}_D, \underline{X}_H)$

- The highest [lowest] fixed point is the stable allocation most preferred by doctors [hospitals] (Theorem 4)

- Starting at  $(X, \emptyset)$  is equivalent to doctor-proposing DA
- Starting at  $(\emptyset, X)$  is equivalent to hospital-proposing DA

- To see equivalence with doctor-proposing DA:

- Define set  $X_H(t)$  as the cumulative contracts offered by doctors to hospitals up to iteration  $t$
- Define  $X_D(t)$  as the cumulative set of contracts that have not yet been rejected by hospitals up to iteration  $t$
- $X_H(t) \cap X_D(t)$  are contracts conditionally accepted
- Using the definition of  $F$ , we have that:

$$\begin{aligned} X_D(t) &= X \setminus R_H(X_H(t-1)), \\ X_H(t) &= X \setminus R_D(X_D(t)). \end{aligned} \tag{1}$$

- If  $|H| \geq 2$ , and  $R_h$  is not isotone for some  $h$ , then there exists a preference profile for doctors and a preference profile for another hospital  $h'$ , which has a single job opening, such that no stable match exists (Theorem 5)

- i.e., substitutes assumption almost needed for existence

### Properties of Generalized DA

- Vacancy chain dynamics can be described by adjusting  $F$  to reject all offers made to a retiring doctor
  - After a doctor retires, a new stable allocation is achieved by iteration on the adjusted  $F$ ; unretired doctors weakly better off and hospitals weakly worse off (Theorem 6)
- The preferences of hospital  $h \in H$  satisfy the law of aggregate demand if for all  $Y \subset Z$ ,  $|C_h(Y)| \leq |C_h(Z)|$ 
  - If more contracts are offered, each hospital should demand weakly more contracts
  - This can be characterized nicely in the matching with wages setting:

**Theorem 7.** *If hospital  $h$ 's preferences are quasi-linear and satisfy the substitutes condition, then they satisfy the law of aggregate demand.*

- A rural hospital theorem also holds in matching with contracts

**Theorem 8.** *If hospital preferences satisfy substitutes and the law of aggregate demand, then for every stable allocation  $(X_D, X_H)$  and every  $d \in D$  and  $h \in H$ ,  $|C_d(X_D)| = |C_d(\bar{X}_D)|$  and  $|C_h(X_H)| = |C_h(\bar{X}_H)|$ .*

- Hospitals and doctors always get the same number of contracts
  - Necessity of law of aggregate demand for this result:

**Theorem 9.** *If there exists a hospital  $h$ , sets  $Y \subset Z \subset X$  such that  $|C_h(Y)| > |C_h(Z)|$ , and  $|H| > 1$ , then there exist singleton preferences for the other hospitals and doctors such that the number of doctors employed by  $h$  is different for two stable matches.*

- Truth-telling is optimal in a few senses
- Let hospitals' preferences satisfy substitutes and the law of aggregate demand, and let the matching algorithm produce the doctor-optimal match. Then, fixing the preferences of the other doctors and of all the hospitals, let  $x$  be the contract that  $d$  obtains by submitting the set of preferences  $P_d : z \succ x$
- Let hospital  $h$  have preferences such that  $|C_h(X)| > |C_h(X \cup \{x\})|$  and let there exist two contracts  $y, z$  such that  $y_D \neq x_D \neq z_D$  and  $y, z \in R_h(X \cup \{x\}) - R_h(X)$ . Then if another hospital  $h'$  exists, there exist singleton preferences for the hospitals besides  $h$  and preferences for the doctors such that it is not a dominant strategy for all doctors to reveal their preferences truthfully.

### Relationship to Proxy Auctions

- The algorithm described by 1 may not converge if substitutes assumption fails, even if a fixed point exists
- The following algorithm is yet another characterization of DA and is suitable for the non-substitutes case, when there is just one hospital (the auctoneer)

$$\begin{aligned} X_D(t) &= X \setminus R_H(X_H(t-1)), \\ X_H(t) &= X_H(t-1) \cup C_D(X_D(t)). \end{aligned} \quad (2)$$

**Theorem 15.** *Under the substitute assumption, with  $X_D(0) = X$  and  $X_H(0) = \emptyset$ , the sequences of pairs  $\{(X_D(t), X_H(t))\}$  generated by the two laws of motion 1 and 2 are identical.*

- The difference between the 1 and 2 is the second equation
  - In 2,  $X_H(t)$  monotonic even if  $R_h$  is not isotone

**Theorem 16.** *When the doctor-offering cumulative offer process with a single hospital terminates at time  $t$  with outcome  $(X_D(t), X_H(t))$ , the hospital's choice  $C_H(X_H(t))$  is a stable collection of contracts.*

### April 24, 2013 Applied Theory Working Group Discussion

Attendees: Sylvain Chassang, Stephen Morris, Ben Brooks, Konstantinos Kalfarentzos, Kai Steverson, Doron Ravid, Olivier Darmouni, Dan Zeltzer, Nemanja Antić